

18/09/18.

LATTICES AND BOOLEAN ALGEBRA.

Partial order relation:-

A relation R on a set A is said to be a partial order relation. If R is reflexive, anti symmetric and transitive.

Result:

Reflexive:- $xRx, \forall x \in R$

Anti symmetric:- $xRy \neq yRx \Rightarrow x=y, \forall x, y \in R.$

Transitive:- $xRy \& yRz \Rightarrow xRz, \forall x, y, z \in R.$

Poset (Partially Ordered set)

A set P together with a partial order relation \leq is called partially Ordered set or POSET and it is denoted by (P, \leq)

Comparable property:-

In a poset for any two elements a, b either $a \leq b$ or $b \leq a$ is called comparable property.

Totally ordered set: (or) Linearly Ordered set: (or) chain

A partially Ordered set (P, \leq) is said to be totally Ordered set. If any 2 elements are comparable.

(i.e.,)

Given any 2 elements ~~except~~ x and y of a poset either

$x \leq y$ or $y \leq x$

Hasse diagram: gram:-

pictorial representation of a poset. is called a Hasse diagram.

1. Draw a Hasse diagram for $(\mathcal{P}(A), \subseteq)$ where

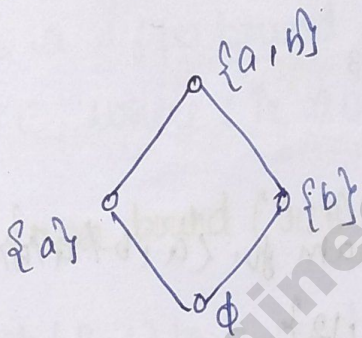
i) $A = \{a, b\}$

ii) $A = \{a, b, c\}$

ans: Hasse diagram for $(\mathcal{P}(A), \subseteq)$

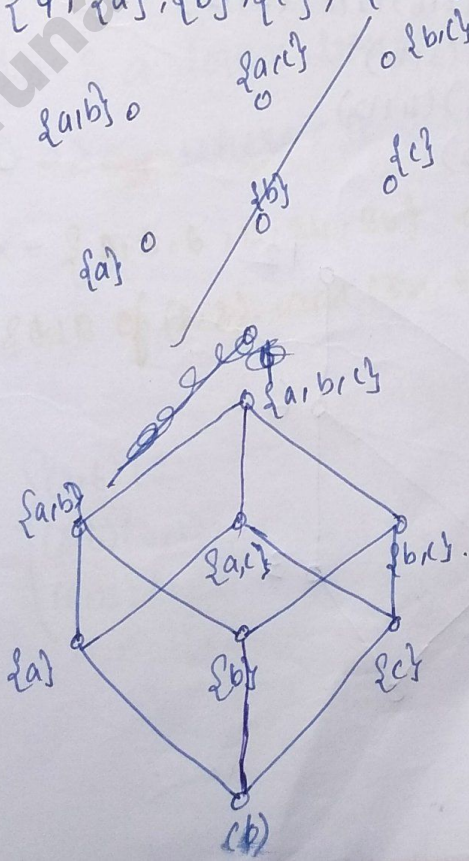
i) $A = \{a, b\}$

$R = \{\emptyset, \{a\}, \{b\}, \{a, b\}, \phi\}$



ii) $A = \{a, b, c\}$

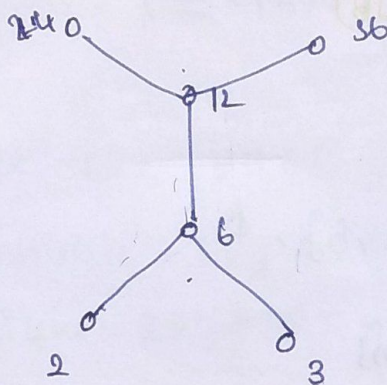
$R = \{\emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{b, c\}, \{a, c\}, \{a, b, c\}\}$



Q. $X = \{2, 3, 6, 12, 24, 36\}$ and the relation R defined on X by $R = \{ \langle a, b \rangle / a|b \}$ Draw Hasse diagram for X, R .

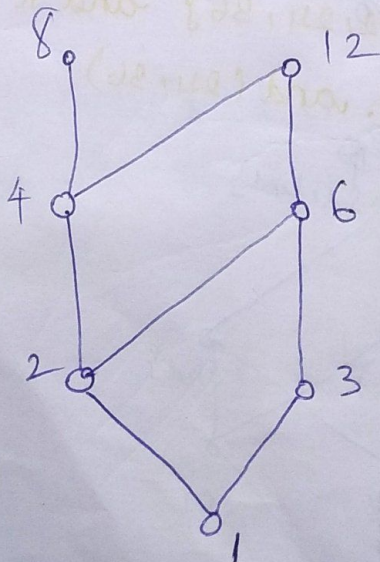
sol.

$$R = \left\{ \begin{array}{l} (2,6) (2,12) (2,24) (2,36) \\ (3,6) (3,12) (3,24) (3,36) \\ (6,12) (6,24) (6,36) \\ (12,24) (12,36) \end{array} \right\}$$



Draw the Hasse diagram for $\langle a, b / a|b \rangle$ on $\{1, 2, 3, 4, 6, 8, 12\}$

$$R = \left\{ \begin{array}{l} (1,2) (1,3) (1,4) (1,6) (1,8) (1,12) \\ (2,4) (2,6) (2,8) (2,12) \\ (3,6) (3,12) \\ (4,8) (4,12) \\ (6,12) \end{array} \right\}$$



5-1
1st
5-8
1
2, 3, 4

Upper bound & Lower bound :-

Let (P, \leq) be a poset and A be any non empty subset of P .

An element $a \in P$ is an upperbound of A .

If $a \geq x \forall x \in A$.

An element $b \in P$ is said to be lowerbound for A if $b \leq x, \forall x \in A$.

Least upper bound :- (LUB)

Let (P, \leq) be a poset and $A \subseteq P$. An element $a \in P$ is said to be least upper bound (LUB) of A if

i) a is an upper bound of A .

ii) $a \leq c$, where c is any other upper bound of A .

Greatest Lower bound (G.L.B) :-

Let (P, \leq) be a poset and $A \subseteq P$. An element $a \in P$ is said to be greatest lower bound (G.L.B) of A if

i) a is a lower bound of A .

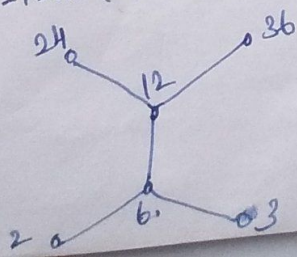
ii) $a \geq c$, where c is any other lower bound of A .

Consider $X = \{2, 3, 6, 12, 24, 36\}$ and $R = \{(a, b) \mid a \mid b\}$

Find LUB & G.L.B of $(2, 3)$ and $(24, 36)$

sol:

$$R = \left\{ \begin{array}{l} (2, 6) (2, 12) (2, 24) (2, 36) \\ (3, 6) (3, 12) (3, 24) (3, 36) \\ (6, 12) (6, 24) (6, 36) \\ (12, 24) (12, 36) \end{array} \right\}$$



$$1. \text{UB } \{2, 3\} = \{6, 12, 24, 36\}$$

$$\text{LUB } \{2, 3\} = \{6\}$$

$$2) \text{UB } \{24, 36\} = \text{does not exist}$$

$$\text{LUB } \{24, 36\} = \text{does not exist}$$

$$3) \text{LB } \{2, 3\} = \text{does not exist}$$

$$\text{GLB } \{2, 3\} = \text{does not exist}$$

$$4) \text{LB } \{24, 36\} = \{2, 3, 6, 12\}$$

$$\text{GLB } \{24, 36\} = \{12\}$$

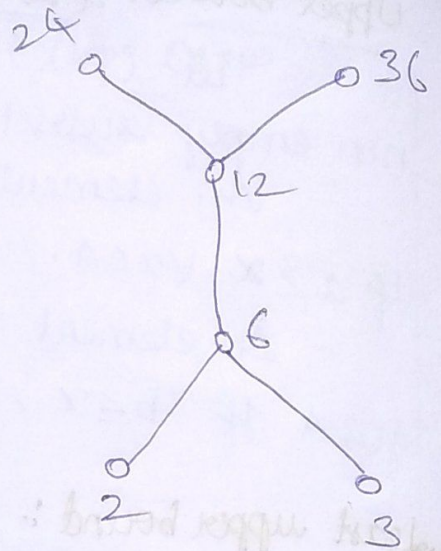
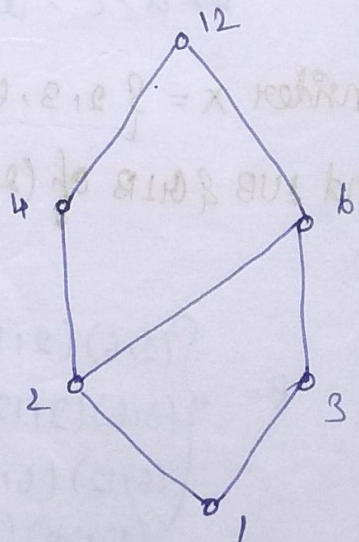
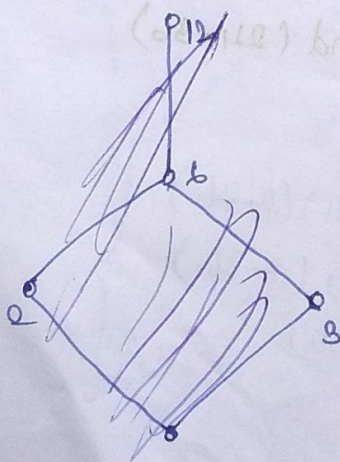
$$\text{GLB } \{24, 36\} = \{12\}$$

Consider $X = \{1, 2, 3, 4, 6, 12\}$ and $R = \{(a, b) \mid a|b\}$ find LUB

& GLB for $\{1, 3\}$ & $\{1, 2, 3\}$, $\{2, 3\}$, $\{2, 3, 6\}$.

sol:

$$R = \left\{ \begin{array}{l} (1, 2) (1, 3) (1, 4) (1, 6) (1, 12) \\ (2, 4) (2, 6) (2, 12) \\ (3, 6) (3, 12) \\ (4, 12) \\ (6, 12) \end{array} \right\}$$



$$1) \text{UB} \{1, 3\} = \{3, 6, 12\}$$

$$\text{LUB} \{1, 3\} = 3$$

$$2) \text{UB} \{1, 2, 3\} = \{3, 6, 12\}$$

$$\text{LUB} \{1, 2, 3\} = 6$$

$$3) \text{UB} \{2, 3\} = \{6, 12\}$$

$$\text{LUB} \{2, 3\} = 6$$

$$4) \text{LB} \{1, 3\} = 1$$

$$\text{GLB} \{1, 3\} = 1$$

$$5) \text{LB} \{1, 2, 3\} = 1$$

$$\text{GLB} \{1, 2, 3\} = 1$$

$$6) \text{LB} \{2, 3\} = 1$$

$$\text{GLB} \{2, 3\} = 1$$

$$7) \text{UB} \{2, 3, 6\} = \{6, 12\}$$

$$\text{LUB} \{2, 3, 6\} = 6$$

$$8) \text{LB} \{2, 3, 6\} = \{2, 3, 6\}$$

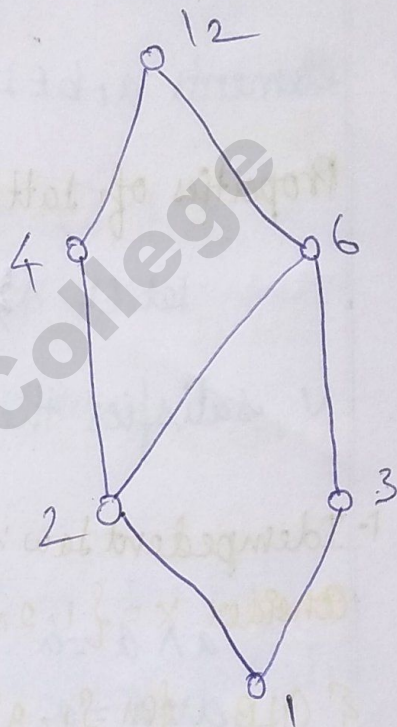
$$\text{GLB} \{2, 3, 6\} = 3$$

$$\text{UB} \{4, 6\} = 12$$

$$\text{LUB} \{4, 6\} = 12$$

$$\text{LB} \{4, 6\} = \{1, 2\}$$

$$\text{GLB} = \{2\}$$



Note:-

$$\text{GIB } (a, b) = a * b \text{ (or) } a \wedge b \text{ (or) } a \cdot b \text{ [meet]}$$

$$\text{LUB } (a, b) = a \oplus b \text{ (or) } a \vee b \text{ (or) } a + b \text{ [join]}$$

Lattice:-

A lattice is a poset L, \leq in which for every pair of elements $a, b \in L$ both GIB and LUB exists.

Properties of Lattice:-

Let (L, \wedge, \vee) be a given lattice then \wedge and \vee satisfies the following conditions, $\forall a, b, c \in L$.

1. Idempotent Law:-

$$a \wedge a = a$$

$$a \vee a = a$$

2. Commutative Law:-

$$a \wedge b = b \wedge a$$

$$a \vee b = b \vee a$$

3. Associative Law:-

$$(a \wedge b) \wedge c = a \wedge (b \wedge c)$$

$$(a \vee b) \vee c = a \vee (b \vee c)$$

4. Absorption Law:-

$$a \wedge (a \vee b) = a$$

$$a \vee (a \wedge b) = a$$

5) $a \vee b = b$ iff $a \leq b$

$a \wedge b = a$ iff $a \leq b$

6) Consistency Law:-

$$a \leq b \Leftrightarrow a \wedge b = a \Leftrightarrow a \vee b = b.$$

Note:- 1) (S_{24}, D) is a lattice

2) (S_{30}, D) is a lattice

Meaning of $a \vee b$ and $a \wedge b$:-

i) $a \leq a \vee b$ and $b \leq a \vee b$

$\therefore a \vee b$ is an upperbound of a and b .

ii) If $a \leq c$ and $b \leq c$ then $a \vee b \leq c$

$\therefore a \vee b$ is the LUB of a and b .

iii) $a \wedge b \leq a$ and $a \wedge b \leq b$

$\therefore a \wedge b$ is lowerbound of a and b .

iv) If $c \leq a$ and $c \leq b$ then $c \leq a \wedge b$.

$\therefore a \wedge b$ is the GIB of a and b .

Meaning of Dual lattice:-

If (L, \leq) is a lattice then the lattice defined by (L, \geq) where the partial ordering \geq is the dual of the partial ordering \leq is called dual lattice.

(ie.,) \wedge by \vee , \vee by \wedge . is called dual statement.

~~Q2~~ Theorem:-

Let (L, \leq) be a distributive lattice then $a \vee b = a \vee c$ and $a \wedge b = a \wedge c \Rightarrow b = c$. $\forall a, b, c \in L$.

Q2)

Cancellation property

Prove that $a \oplus b = a \oplus c$ and $a * b = a * c \Rightarrow b = c$.

Proof:-

W.K.T

$$b = b \vee b$$

$$= b \vee (a \wedge b)$$

$$= b \vee (a \wedge c) \quad [\because a \wedge b = a \wedge c]$$

$$= (b \vee a) \wedge (b \vee c)$$

$$\begin{aligned}
 &= \varnothing \\
 &= (a \vee c) \wedge (b \vee c) \quad [\because a \vee b = a \vee c] \\
 &= (a \wedge b) \vee c \\
 &= (a \wedge c) \vee c = c \quad [\because \text{Absorption law}]
 \end{aligned}$$

$$b = c$$

2) Theorem:-

State and prove distributive inequalities

Statement:-

Let (L, \leq) be a lattice. For any $a, b, c \in L$, then

$$a \wedge (b \vee c) \geq (a \wedge b) \vee (a \wedge c)$$

$$a \vee (b \wedge c) \leq (a \vee b) \wedge (a \vee c)$$

Proof:-

W.K.T $a \wedge b \leq a$ and $a \wedge b \leq b \leq b \vee c$

$\therefore a \wedge b$ is a lower bound of a and $b \vee c$

$$\Rightarrow (a \wedge b) \leq a \wedge (b \vee c) \text{ --- (1)}$$

Again $a \wedge c \leq a$ and $a \wedge c \leq c \leq b \vee c$

$\therefore a \wedge c$ is a L.B of a and $b \vee c$

$$\Rightarrow a \wedge c \leq a \wedge (b \vee c) \text{ --- (2)}$$

from (1) & (2)

$a \wedge (b \vee c)$ is an UB of $a \wedge b$ and $a \wedge c$.

But $(a \wedge b) \vee (a \wedge c)$ is the LUB of $a \wedge b$ and $a \wedge c$.

$$\Rightarrow a \wedge (b \vee c) \geq (a \wedge b) \vee (a \wedge c)$$

By applying Duality,

$$a \vee (b \wedge c) \leq (a \vee b) \wedge (a \vee c)$$

3) show that a chain is a lattice

Proof:-

Let (L, \leq) be a chain

If $a, b \in L$, $a \leq b$ (or) $b \leq a$

If $a \leq b$, $a \wedge b = a$, $a \vee b = b$

\therefore LUB and GLB of a and b exists

If $b \leq a$, $b \wedge a = b$, $b \vee a = a$

\therefore LUB and GLB of a and b exists

Hence every pair of elements has a GLB and LUB.

\therefore Every chain is a lattice.

Distributive lattice:-

$$a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$$

$$a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$$

Lattice Homomorphism:-

Let (L_1^*, \oplus) and S, \wedge, \vee with be two lattices a mapping $g: L \rightarrow S$ is called lattice

homomorphism. If $g(a * b) = g(a) \wedge g(b)$

$$g(a \oplus b) = g(a) \vee g(b)$$

Modular lattice:-

The lattice L is said to be a modular

lattice if $\forall a, b, c \in L$.

$$a \leq c \Rightarrow a \vee (b \wedge c) = (a \vee b) \wedge c$$

1) Theorem :-

Prove that every distributive lattice is modular lattice.

Proof :-

Let (L, \leq) be a distributive lattice.

Let $a, b, c \in L$ such that $a \leq c$.

To prove: L is modular lattice.

i.e. To prove: $a \vee (b \wedge c) = (a \vee b) \wedge c$.

Now $a \leq c \Rightarrow a \vee c = c$.

$$\begin{aligned} \text{L.H.S } a \vee (b \wedge c) &= (a \vee b) \wedge (a \vee c) \\ &= (a \vee b) \wedge c \\ &= \text{R.H.S.} \end{aligned}$$

$\therefore L$ is a modular lattice.

2) Theorem :-

Show that every chain is a distributive lattice.

Proof :-

Let (L, \leq) be a chain.

If $a, b, c \in L$, then $a \leq b$ or $b \leq a$.

i) Suppose that

$a \leq b$ (or) $a \leq c$, then $a \leq b \vee c$

$$\begin{aligned} \text{Let } a \wedge (b \vee c) &= a \wedge a \\ &= a \end{aligned}$$

$$\therefore a \leq b \Rightarrow a \wedge b = a$$

$$a \leq c \Rightarrow a \wedge c = a.$$

$$a \leq b \vee c \Rightarrow a \wedge (b \vee c)$$

R.H.S

ii) Suppose that

~~$b \leq a$ (or) $c \leq a$~~

$$(a \wedge b) \vee (a \wedge c)$$

$$= a \vee a$$

$$= a$$

$$\text{L.H.S} = \text{R.H.S}$$

By duality, $a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$

ii) Suppose that

$$b \leq a \text{ (or) } c \leq a$$

$$\text{then } b \vee c \leq a$$

Now

$$b \leq a \Rightarrow b \wedge a = b$$

$$c \leq a \Rightarrow c \wedge a = c.$$

$$b \vee c \leq a \Rightarrow (b \vee c) \wedge a = b \vee c$$

$$\text{L.H.S } a \wedge (b \vee c) = b \vee c$$

$$\text{RHS } (a \wedge b) \vee (a \wedge c) = b \vee c$$

$$\Rightarrow \text{L.H.S} = \text{RHS}$$

$$\therefore a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$$

By dualizing

$$a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$$

Complement of an element:-

Let L be a bounded lattice with $LB = 0$

and $UB = 1$

$$\therefore LB = 1 \text{ \& } UB = 0.$$

The element $x \in L$ is called the complement

of $a \in L$.

$$\text{iff } a \wedge x = 0 \text{ and } a \vee x = 1.$$

Complemented lattice:-

A lattice L is said to be complemented if it is bounded and every element in it has at least 1 complement.

Bounded Lattice:-

A lattice L is said to be bounded if it has both LB and UB, If L is a bounded lattice then,

$$a \vee 0 = a, a \wedge 0 = 0$$

$$a \vee 1 = 1, a \wedge 1 = a.$$

Theorem:-

prove that in a bounded distributive lattice the complement of any element is unique.

proof:-

Let L be a bounded distributive lattice.

Let b & c be complements of an element $a \in L$.

To prove $b = c$.

Since b & c are complements of 'a' we have.

$$a \wedge b = 0 \rightarrow \textcircled{1}, a \vee b = 1 \rightarrow \textcircled{2}.$$

$$a \wedge c = 0 \rightarrow \textcircled{3}, a \vee c = 1 \rightarrow \textcircled{4}.$$

Now,

$$b = b \wedge 1$$

$$= b \wedge (a \vee c) \quad \because \text{by } \textcircled{4}$$

$$= (b \wedge a) \vee (b \wedge c)$$

$$= 0 \vee (b \wedge c) \quad \because \text{by } \textcircled{1}.$$

$$= (a \wedge c) \vee (b \wedge c) \quad \because \text{by } \textcircled{3}$$

$$= (a \vee b) \wedge c \quad \because \text{by } \textcircled{2}.$$

$$= 1 \wedge c = c$$

$$\boxed{\therefore b = c}$$

\therefore The complement of any element is unique.

De Morgan's Law:-

Prove De Morgan's law in a bounded, complemented and distributive lattice.

Sol:

De Morgan's laws are given by

$$(a \vee b)' = a' \wedge b' \quad (\text{or}) \quad (a \oplus b)' = a' * b'$$

$$(a \wedge b)' = a' \vee b' \quad (\text{or}) \quad (a * b)' = a' \oplus b'$$

To prove: $(a \vee b)' = a' \wedge b'$

w.k.t $a \vee a' = 1 \rightarrow \textcircled{1}$. $a \wedge a' = 0 \rightarrow \textcircled{2}$

$b \vee b' = 1 \rightarrow \textcircled{3}$. $b \wedge b' = 0 \rightarrow \textcircled{4}$

and $1 \vee x = 1 \rightarrow \textcircled{5}$. $0 \wedge x = 0 \rightarrow \textcircled{6}$

consider, $(a \vee b) \vee (a' \wedge b')$
 $= (a \vee b \vee a') \wedge (a \vee b \vee b')$
 $= (1 \vee b) \wedge (a \vee 1) \quad \dots \text{by } \textcircled{1} \& \textcircled{3}$
 $= 1 \wedge 1 \quad \dots \text{by } \textcircled{5}$

$(a \vee b) \vee (a' \wedge b') = 1 \rightarrow \textcircled{7}$

Now $(a \vee b) \wedge (a' \wedge b')$

$= (a \wedge a' \wedge b') \vee (b \wedge a' \wedge b')$

$= (0 \wedge b') \vee (0 \wedge a') \quad \dots \text{by } \textcircled{2} \& \textcircled{4}$

$= 0 \vee 0 \quad \dots \text{by } \textcircled{6}$

$= 0 \rightarrow \textcircled{8}$

from $\textcircled{7} \& \textcircled{8}$

$\therefore a' \wedge b'$ is the complement of $a \vee b$

Hence $(a \vee b)' = a' \wedge b'$

By duality $(a \wedge b)' = a' \vee b'$

Boolean algebra:

A complemented distributive lattice is called a boolean algebra.

Abstract definition:-

Let B be a non-empty set that contains two special elements 0 (element) and 1 (unity) and on which we defined closed binary operations $+$, \cdot , and a unary operation $' - '$ (complement) then $(B, +, \cdot, 0, 1)$ is called a boolean algebra.

If the following conditions are satisfied

For $x, y, z \in B$

1. $x + y = y + x, x \cdot y = y \cdot x$.

2. $x + (y + z) = (x + y) + z, x \cdot (y \cdot z) = (x \cdot y) \cdot z$.

3. $x + 0 = x, x \cdot 1 = x$.

4. $x + \bar{x} = 1, x \cdot \bar{x} = 0$.

5. $x + (y \cdot z) = (x + y) \cdot (x + z)$

$x \cdot (y + z) = (x \cdot y) + (x \cdot z)$

Laws of Boolean Algebra:

1. $x \cdot 0 = 0; x + 1 = 1$ [Dominance law]

2. $x \cdot (x + y) = x; x + (x \cdot y) = x$.

3. $x \cdot y = z$
 $\bar{x} \cdot y = \bar{x} \cdot z$ } $\Rightarrow y = z$.

$x + y = x + z$
 $\bar{x} + y = \bar{x} + z$ } $\Rightarrow y = z$

cancellation law.

4. $x \cdot (y \cdot z) = (x \cdot y) \cdot z$.

$x + (y + z) = (x + y) + z$.

5. $x + y = 1; x \cdot y = 0$
 $\Rightarrow x = \bar{y}$ (or) $y = \bar{x}$

$$v) x+y=1; x \cdot y=0$$

$$x = \bar{y} \text{ (or) } y = \bar{x}$$

$$vi) \bar{\bar{x}} = x.$$

$$vii) (\overline{x+y}) = \bar{x} \cdot \bar{y}$$

$$(\overline{x \cdot y}) = \bar{x} + \bar{y}$$

$$viii) \bar{0} = 1; \bar{1} = 0.$$

show that demorgan's law is valid in boolean algebra.

sol:

Demorgan's law are $(a+b)' = a' \cdot b'$

$$(a \cdot b)' = a' + b'$$

i) consider $(a+b) + (a' \cdot b')$

$$= (a+b+a') \cdot (a+b+b')$$

$$= (1+b) \cdot (a+1)$$

$$= 1 \cdot 1$$

$$= 1.$$

$$(a+b) \cdot (a' \cdot b') = (a \cdot a' \cdot b') + (b \cdot a' \cdot b')$$

$$= (0 \cdot b') + (a' \cdot 0)$$

$$= 0 + 0$$

$$= 0.$$

$\therefore a' \cdot b'$ is the complement of $(a+b)$

$$\text{Hence } (a+b)' = a' \cdot b'.$$

ii) $(a \cdot b) + (a' + b')$

$$= (a+a'+b') \cdot (b+a'+b')$$

$$= (1+b') \cdot (a'+1)$$

$$= 1 \cdot 1$$

$$= 1.$$

$$(a \cdot b) \cdot (a' + b') = (a \cdot b \cdot a') + (a \cdot b \cdot b')$$

$$= (0 \cdot b) + (a \cdot 0)$$

$$= 0 + 0$$

$$= 0.$$

$\therefore a' + b'$ is the complement of $a \cdot b$

$$\text{Hence } (a \cdot b)' = a' + b'.$$

Theorem:-

If any boolean algebra, show that $(a+b')(b+c')(c+a)$

$$(b+c')(c+a') = (a'+b)(b'+c)(c'+a)$$

Sol.

LHS

$$(a+b')(b+c')(c+a')$$

$$= (a+b')(bc + ba' + c'c + c'a')$$

$$= abc + aba' + ac'c + ac'a' + b'bc +$$

$$b'ba' + b'c'c + b'c'a'$$

$$= abc + 0 + 0 + 0 + 0 + 0 + 0 + 0 + b'c'a'$$

$$= abc + b'c'a' \text{ --- (1)}$$

RHS

$$(a'+b)(b'+c)(c'+a)$$

$$= (a'+b)(b'c' + b'a + cc' + ca)$$

$$= a'b'c' + a'b'a + a'cc' + a'ca + bb'c' + bb'a + bcc' + bca$$

$$= a'b'c' + 0 + 0 + 0 + 0 + 0 + 0 + bca$$

$$= abc + a'b'c' \text{ --- (2)}$$

from (1) & (2)

$$L.H.S = R.H.S.$$

$$(a+b')(b+c')(c+a') = (a'+b)(b'+c)(c'+a)$$

$$a \neq b \Rightarrow a \neq b = 0 \neq 0$$

Show that in a boolean algebra,
 $a \leq b \Leftrightarrow a \wedge \bar{b} = 0 \Leftrightarrow \bar{a} \vee b = 1 \Leftrightarrow \bar{b} \leq \bar{a}$.

Sol:

i) To prove:

$$a \leq b \Rightarrow a \wedge \bar{b} = 0.$$

w.k.t $a \leq b \Rightarrow a \wedge b = a \rightarrow$ (i) $a \vee b = b \rightarrow$ (ii).

$$a \wedge \bar{b} = a \wedge (b \wedge \bar{b}) \quad \therefore \text{by (i)}$$

$$= a \wedge 0$$

$$= 0$$

$$\therefore a \wedge \bar{b} = 0$$

ii) To prove.

$$a \wedge \bar{b} = 0 \Rightarrow \bar{a} \vee b = 1.$$

Now $a \wedge \bar{b} = 0$.

Taking complement on b/s.

$$\overline{a \wedge \bar{b}} = \bar{0}$$

$$\therefore \bar{a} \vee b = 1.$$

iii) To prove: $\bar{a} \vee b = 1 \Rightarrow \bar{b} \leq \bar{a}$

Now, $\bar{a} \vee b = 1$

$$(\bar{a} \vee b) \wedge \bar{b} = 1 \wedge \bar{b}$$

$$(\bar{a} \wedge \bar{b}) \vee (b \wedge \bar{b}) = 1 \wedge \bar{b}$$

$$(\bar{a} \wedge \bar{b}) \vee 0 = 1 \wedge \bar{b}$$

$$\bar{a} \wedge \bar{b} = \bar{b}$$

$$\bar{b} \leq \bar{a}$$

iv) To prove:

$$\bar{b} \leq \bar{a} \Rightarrow a \leq b$$

$$\bar{b} \leq \bar{a} \Rightarrow \bar{a} \wedge \bar{b} = \bar{b}$$

$$\Rightarrow \overline{\bar{a} \wedge \bar{b}} = \bar{\bar{b}}$$

$$\Rightarrow a \vee b = b$$

$$\Rightarrow a \leq b.$$

In a boolean algebra show that the following statements are equivalent. For any a and b .

i) $a+b=b$ ii) $a \cdot b = a$ iii) $a'+b=1$ iv) $a \cdot b' = 0$

v) $a \leq b$

Proof:

i) \Rightarrow ii):

Assume that $a+b=b$.

$$a = a \cdot (a+b)$$

$$\boxed{= a \cdot b = a}$$

ii) \Rightarrow iii):

Assume that $a \cdot b = a$

$$a'+b = (a \cdot b)'+b$$

$$= a'+b'+b$$

$$= a'+1$$

$$\boxed{a'+b = 1}$$

iii) \Rightarrow iv):

Assume that $a'+b=1$

$$(a'+b)' = (1)' = 0$$

$$a \cdot b' = 0$$

iv) \Rightarrow v):

Assume that $a \leq b$

$$a+b = a \cdot b + 1 \cdot b$$

$$= a \cdot b + 1 \cdot b$$

$$= (a+1) \cdot b$$

$$= 1 \cdot b$$

$$= b$$

v) \Rightarrow i): Assume that $a \cdot b' = 0$

$$\text{then } a \cdot b = a \cdot b + 0$$

$$= a \cdot b + a \cdot b'$$

$$= a \cdot (b+b')$$

$$= a \cdot 1$$

$$a \cdot b = a$$

$$= a \leq b$$